

CHAPTER 1

QUANTIFIERS & PROOF TECHNIQUE [Part 4]

QUANTIFIERS

- Most of the statements in mathematics and computer science are not described properly by the propositions.
- Since most of the statements in mathematics and computer science use **variables**, the system of logic must be extended to include statements with the variables.

QUANTIFIERS (cont'd)

- Let $P(x)$ is a statement with variable x and $\mathbf{A}$ is a set.
- P a **propositional function** or also known as **predicate** if for each x in $\mathbf{A}$, $P(x)$ is a proposition.
- Set $\mathbf{A}$ is the **domain of discourse** of P .
- Domain of discourse: the particular domain of the variable in a propositional function.

QUANTIFIERS (cont'd)

A **predicate** is a statement that contains **variables**.

Example:

$$P(x): x > 3$$

$$Q(x, y): x = y + 3$$

$$R(x, y, z): x + y = z$$

Example:

The following are **propositional functions**:

1) $x^2 + 4x$ is an odd integer

(domain of discourse is set of positive numbers).

2) $x^2 - x - 6 = 0$

(domain of discourse is set of real numbers).

3) The university rated as Research University in Malaysia

(domain of discourse is set of research university in Malaysia).

QUANTIFIERS (cont'd)

A predicate becomes a proposition if the variable(s) contained is(are)

- Assigned specific value(s)
- Quantified

Example:

i) $P(x): x > 3$

What are the truth values of $P(4)$ and $P(2)$?

ii) $Q(x,y): x = y + 3$

What are the truth values of $Q(1,2)$ and $Q(3,0)$?

QUANTIFIERS (cont'd)

Two types of quantifiers:

a) Universal

b) Existential

Universal Quantifier

- Let **A** be a propositional function with domain of discourse **B**. The statement:

for every x , $A(x)$

is **universally quantified statement**

- Symbol $\forall$ called a **universal quantifier** is used “**for every**”.
- Can be read as “*for all*”, “*for any*”.

Universal Quantifier (cont'd)

- The statement can be written as

$$\forall x \mathbf{A}(x)$$

- Above statement is **true** if $\mathbf{A}(x)$ is true **for every** x in $\mathbf{B}$ and it **false** if $\mathbf{A}(x)$ is false for **at least one** x in $\mathbf{B}$.
- A value x in the domain of discourse that makes the statement $\mathbf{A}(x)$ false is called as a **counterexample** to the statement.

Example:

Let the universally quantified statement is

$$\forall x (x^2 \geq 0)$$

Domain of discourse is the set of real numbers.
Determine the truth value of this statement.

Answer:

This statement is **true** because **for every** real number x , it is true that the square of x is positive or zero.

Example:

Let the universally quantified statement is

$$\forall x (x^2 \leq 9)$$

Domain of discourse is a set $B = \{1, 2, 3, 4\}$

Determine the truth value for this statement.

Answer:

This statement is **false** and the counterexample is 4, because when $x = 4$, the statement produce false value.

Universal Quantifier (cont'd)

- Easy to prove a universally quantified statement is true or false if the domain of discourse is not too large.
- What happen if the domain of discourse contains a large number of elements?
- For example, a set of integer from 1 to 100, the set of positive integers, the set of real numbers or a set of students in Faculty of Computing. It will be hard to show that every element in the set is true.
- Therefore, we need to use **existential quantifier**.

Existential Quantifiers (cont'd)

- Let **A** be a propositional function with domain of discourse **B**. The statement

There exist x , $A(x)$

is **existentially quantified statement**

- Symbol $\exists$ called an **existential quantifier** is used “**there exist**”.
- Can be read as “*for some*”, “*for at least one*”.

Existential Quantifiers (cont'd)

- The statement can be written as

$$\exists x \mathbf{A}(x)$$

- Above statement is **true** if it possible to find **at least one** x in **B** that makes $\mathbf{A}(x)$ true, and it will be **false** if **every** x in **B** that makes the statement $\mathbf{A}(x)$ false.
- Just find one x that makes $\mathbf{A}(x)$ true!

Example:

Let the existentially quantified statement is

$$\exists x \left(\frac{x}{x^2 + 1} = \frac{2}{5} \right)$$

Domain of discourse is the set of real numbers.

Determine the truth value for this statement.

Answer:

Statement is **true** because it is possible to find at least one real number x to make the proposition true.

For example, if $x = 2$, we obtain the true proposition as below.

$$\left(\frac{x}{x^2 + 1} = \frac{2}{5} \right) = \left(\frac{2}{2^2 + 1} = \frac{2}{5} \right)$$

Negation of Quantifiers

Distributing a negation operator across a quantifier changes a universal to an existential and vice versa.

$$\neg (\forall x P(x)) \Rightarrow \exists x \neg P(x)$$

$$\neg (\exists x P(x)) \Rightarrow \forall x \neg P(x)$$

Example:

Let $P(x) = x$ is taking Discrete Structure course with the domain of discourse is the set of all students.

- $\forall x P(x)$: All students are taking Discrete Structure course.
- $\exists x P(x)$: There is some students who are taking Discrete Structure course.
- $\forall x \neg P(x)$: All students are **not** taking Discrete Structure course.
- $\exists x \neg P(x)$: There are some students who are **not** taking Discrete Structure course.

PROOF TECHNIQUES

Mathematical systems consists:

- ✓ Axioms - assumed to be true.
- ✓ Definitions - used to create new concepts.
- ✓ Undefined terms - some terms that are not explicitly defined.
- ✓ Theorem – proposition that has been proved to be true.

Theorem:

- It can be stated as follows:
 - As facts
 - As implications
 - As bi-implications
- An argument that establishes the truth of a theorem is called a **proof**.

PROOF TECHNIQUES (cont'd)

There are several techniques to proof theorem:

- i. Direct proof
- ii. Indirect proof
- iii. Proof by contradiction

Direct Proof

- Direct proof or proof by direct method assumes that $p(x_1, x_2, \dots, x_n)$ is true and then using $p(x_1, x_2, \dots, x_n)$ as well as other axioms, definitions and theorems, show directly that $q(x_1, x_2, \dots, x_n)$ is true.

Symbol: $\forall x (P(x) \rightarrow Q(x))$

- Procedure for direct proof:
 - ✓ Let **D** is the domain of discourse, select a particular, but arbitrarily chosen, member of the domain **D**.
 - ✓ Then show that the statement $\forall x (P(x) \rightarrow Q(x))$ is true by assuming that $P(x)$ is true and then show that $Q(x)$ is also true.

Example:

Use direct proof to prove the following theorem.

For all integer x , if x is odd, then x^2 is odd.

Solution:

We can verify the theorem by substituting x with certain values.

For example,

$$x = 3 \rightarrow x^2 = 9 \text{ (both odd);}$$

$$x = 511 \rightarrow x^2 = 261121 \text{ (both odd)}$$

However, verifying a given theorem for a particular value is not a proof. Therefore, we must prove that the theorem is true for an arbitrary value. So, let

$$P(x) = x \text{ is an odd integer}$$

$$Q(x) = x^2 \text{ is an odd integer}$$

Solution (cont'd)

Symbolically, $\forall x(P(x) \rightarrow Q(x))$ with domain of discourse is the set of all integers. Let a is an odd integer.

$$\Rightarrow a = 2n + 1 \quad \text{for some integer } n$$

$$\Rightarrow a^2 = (2n + 1)^2$$

$$\Rightarrow a^2 = 4n^2 + 4n + 1$$

$$\Rightarrow a^2 = 2(2n^2 + 2n) + 1$$

$$\Rightarrow a^2 = 2m + 1 \quad \text{where } m = 2n^2 + 2n \text{ is an integer}$$

$$\Rightarrow a^2 \quad \text{is an odd integer}$$

Indirect Proof

Consider the implication $p \rightarrow q$ which is **equivalent** to the implication $\neg q \rightarrow \neg p$.

Therefore, in order to show that $p \rightarrow q$ is true, we can also show that the implication $\neg q \rightarrow \neg p$ is true.

To show that $\neg q \rightarrow \neg p$ is true, assume that the negation of q is true and prove that the negation of p is true.

Example:

Use indirect proof to prove the following:

Let n be an integer, if n^2+3 is odd, then n is even.

Solution:

Let,

$P(n) = n^2+3$ is an odd integer

$Q(n) = n$ is an even integer.

Symbolically, $\forall x(P(x) \rightarrow Q(x))$ with domain of discourse is the set of all integers.

Assume n is a particular integer but arbitrary chosen element from the domain of discourse.

Solution (cont'd)

For this n , suppose $\neg Q(n)$ is true, we need to show that $\neg P(n)$ is true.

Because $\neg Q(n)$ is true, n is not even. Then, n is odd.

So, $n = 2k + 1$ for some integer k .

$$n = 2k + 1$$

$$n^2 = (2k + 1)^2$$

$$n^2 = 4k^2 + 4k + 1$$

$$n^2 + 3 = (4k^2 + 4k + 1) + 3$$

$$n^2 + 3 = 4k^2 + 4k + 4$$

$$n^2 + 3 = 2(2k^2 + 2k + 2)$$

Because k is integer, thus $t = 2k^2 + 2k + 2$ is an integer. So, $n^2 + 3 = 2t$, which is multiple of 2. Therefore, $n^2 + 3$ is even integer, $\neg P(n)$ is true.

Proof by Contradiction

In a proof by contradiction, we assume that the conclusion is not true then arrive at a contradiction.

Example:

Prove the following theorem:

There are infinitely many prime numbers.

Solution:

Assume to the **contrary** that there are **only finite** many prime numbers and all of them are listed as: $p_1, p_2, \dots, p_n$

Consider the number $q = p_1, p_2, \dots, p_{n+1}$. The number q is either prime or not divisible. If we divide any of the listed primes, p_i into q , there would result a remainder of 1 for each $i = 1, 2, \dots, n$. Thus q is not divisible, but not listed above. Therefore, q is a prime.

Contradiction! Therefore, there are infinitely many primes numbers.

Exercise

Let $P(x, y) = (x * y)^2 \geq 1$. Given the domain of discourse for x and y is set of integer, Z .

Determine the truth value of the following statements. Give the value of x and y that make the statement TRUE or FALSE.

$$a) \exists x \exists y P(x, y)$$

$$b) \forall x \forall y P(x, y)$$