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SCSI1013: Discrete Structures

CHAPTER 1

Part 3:
Fundamental and Elements of Logic

2017/2018 – SEM. 1 : nzah@utm.my

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Why Are We Studying Logic?

Some of the reasons:

- Logic is the foundation for computer operation
- Logical conditions are common in programs:
Example:
Selection: if (score <= max) { ... }
Iteration: while (i < limit && list[i] != sentinel) ...
- All manner of structures in computing have properties that need to be proven (and proofs that need to be understood).
Examples: Trees, Graphs, Recursive Algorithms, . . .
- Programs can be proven correct.
- Computational linguistics must represent and reason about human language, and language represents thought (and thus also logic).

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PROPOSITION

A **statement** or a **proposition**, is a declarative sentence that is **either TRUE or FALSE, but not both**.

Example:

- 4 is less than 3.
- 7 is an even integer.
- Washington, DC, is the capital of United State.

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Example

i) Why do we study mathematics?
 ii) Study logic.
 iii) What is your name?
 iv) Quiet, please.

The above sentences are not propositions. Why ?

(i) & (iii) : is question, not a statement.
 (ii) & (iv) : is a command.

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Example

i) The temperature on the surface of the planet Venus is 800 F.
 ii) The sun will come out tomorrow.

Propositions? Why?

i) Is a statement since it is either true or false, but not both.
 ii) However, we do not know at this time to determine whether it is true or false.

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CONJUNCTIONS

Conjunctions are:

- Compound propositions formed in English with the word “**and**”,
- Formed in logic with the caret symbol (“**∧**”), and
- True only when both participating propositions are true.



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CONJUNCTIONS (cont.)

TRUTH TABLE: This tables aid in the evaluation of compound propositions.

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

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Example

p : 2 is an even integer } propositions
 q : 3 is an odd number }

$p \wedge q$ } symbols } statements
 : 2 is an even integer and 3 is an odd number

p : today is Monday ; q : it is hot

$p \wedge q$: today is Monday and it is hot

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Example

Proposition:-
 p : 2 divides 4
 q : 2 divides 6

Symbol: Statement:-
 $p \wedge q$: 2 divides 4 and 2 divides 6.
 or,
 $p \wedge q$: 2 divides both 4 and 6.

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Example


Proposition:-
 p : 5 is an integer
 q : 5 is not an odd integer

Symbol: Statement:-
 $p \wedge q$: 5 is an integer and 5 is not an odd integer.
 or,
 $p \wedge q$: 5 is an integer but 5 is not an odd integer.

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DISJUNCTION

- Compound propositions formed in English with the word "or",
- Formed in logic with the caret symbol (" \vee "), and,
- True when one or both participating propositions are true.



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DISJUNCTION (cont.)

Let p and q be propositions. The **disjunction** of p and q , written $p \vee q$ is the statement formed by putting statements p and q together using the word "or". The symbol \vee is called "or"

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UTM **DISJUNCTION** (cont.)

The truth table for $p \vee q$:

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

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UTM **Example**

i) p : 2 is an integer ; q : 3 is greater than 5
 $p \vee q \rightarrow$ 2 is an integer or 3 is greater than 5

ii) p : $1+1=3$; q : A decade is 10 years
 $p \vee q \rightarrow$ $1+1=3$ or a decade is 10 years

iii) p : 3 is an even integer ; q : 3 is an odd integer
 $p \vee q \rightarrow$ 3 is an even integer or 3 is an odd integer ; or
 3 is an even integer or an odd integer

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UTM **NEGATION**

Negating a proposition simply flips its value. Symbols representing negation include: $\neg x, \bar{x}, \sim x, x'$ (NOT)

Let p be a proposition.
 The negation of p , written $\neg p$ is the statement obtained by negating statement p .

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UTM **NEGATION** (cont.)

The truth table of $\neg p$:

p	$\neg p$
T	F
F	T

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UTM **Example**

p : 2 is positive
 $\neg p \rightarrow$ 2 is not positive.

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UTM **Exercise (1)**

p : It will rain tomorrow ; q : it will snow tomorrow

Give the negation of the following statement and write the symbol.

It will rain tomorrow or it will snow tomorrow.

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Exercise (2)

In each of the following, form the conjunction and the disjunction of p and q by writing the symbol and the statements.

i) p : I will drive my car
 q : I will be late

ii) p : NUM > 10
 q : NUM ≤ 15

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Exercise (3)

Suppose x is a particular real number. Let p , q and r symbolize " $0 < x$ ", " $x < 3$ " and " $x = 3$ ", respectively. Write the following inequalities symbolically:

a) $x ≤ 3$
 b) $0 < x < 3$
 c) $0 < x ≤ 3$

Solution:
 a) $q ∨ r$
 b) $p ∧ q$
 c) $p ∧ (q ∨ r)$

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Exercise (4)

State either TRUE or FALSE if p and r are TRUE and q is FALSE.

a) $∼ p ∧ (q ∨ r)$

a) $(r ∧ ∼ q) ∨ (p ∨ r)$

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CONDITIONAL PROPOSITIONS

Let p and q be propositions.

"if p , then q "

is a statement called a **conditional proposition**, written as

$p → q$

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CONDITIONAL PROPOSITIONS (cont.)

The **truth table** of $p → q$
 => Cause and effect relationship

FALSE if p = True and q = false

p	q	$p → q$
T	T	T
T	F	F
F	T	T
F	F	T

TRUE if both true or p =false for any value of q

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Example

p : today is Sunday ; q : I will go for a walk

$p → q$: If today is Sunday, then I will go for a walk.

p : I get a bonus ; q : I will buy a new car

$p → q$: If I get a bonus, then I will buy a new car

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Example

p : $x/2$ is an integer.
 q : x is an even integer.

$p \rightarrow q$: if $x/2$ is an integer, then x is an even integer.

BICONDITIONAL

Let p and q be propositions.

“ p if and only if q ”

is a statement called a **biconditional proposition**, written as

$p \leftrightarrow q$

BICONDITIONAL (cont.)

The **truth table** of $p \leftrightarrow q$:

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

Example

p : my program will compile
 q : it has no syntax error.

$p \leftrightarrow q$: My program will compile if and only if it has no syntax error.

p : x is divisible by 3
 q : x is divisible by 9

$p \leftrightarrow q$: x is divisible by 3 if and only if x is divisible by 9.

LOGICAL EQUIVALENCE

- The compound propositions Q and R are made up of the propositions p_1, \dots, p_n .
- Q and R are logically equivalent and write, **$Q \equiv R$** provided that given any truth values of p_1, \dots, p_n , either Q and R are **both true** or Q and R are **both false**.

Example

$Q = p \rightarrow q$
 $R = \neg q \rightarrow \neg p$
 Show that, $Q \equiv R$

The **truth table** shows that, $Q \equiv R$

p	q	$p \rightarrow q$	$\neg q \rightarrow \neg p$
T	T	T	T
T	F	F	F
F	T	T	T
F	F	T	T

Example

Show that,
 $\neg(p \rightarrow q) \equiv p \wedge \neg q$

The truth table shows that,
 $\neg(p \rightarrow q) \equiv p \wedge \neg q$

p	q	$\neg(p \rightarrow q)$	$p \wedge \neg q$
T	T	F	F
T	F	T	T
F	T	F	F
F	F	F	F

PRECEDENCE OF LOGICAL CONNECTIVES

Precedence of logical connectives is as follows:

Connective	Symbol	Precedence
not	\neg	Highest
and	\wedge	
or	\vee	
if...then	\rightarrow	
If and only if	\leftrightarrow	Lowest

Example

Construct the truth table for,
 $A = \neg(p \vee q) \rightarrow (q \wedge p)$

Solution:

p	q	$(p \vee q)$	$\neg(p \vee q)$	$(q \wedge p)$	A
T	T	T	F	T	T
T	F	T	F	F	T
F	T	T	F	F	T
F	F	F	T	F	F

Exercise (5)

Construct the truth table for each of the following statements:

- $\neg p \wedge q$
- $\neg(p \vee q) \rightarrow q$
- $\neg(\neg p \wedge q) \vee q$
- $(p \rightarrow q) \rightarrow (\neg q \rightarrow \neg p)$

LOGIC & SET THEORY

Logic and set theory go very well together. The previous definitions can be made very succinct:

- $x \notin A$ if and only if $\neg(x \in A)$
- $A \subseteq B$ if and only if $(x \in A \rightarrow x \in B)$ is True
- $x \in (A \cap B)$ if and only if $(x \in A \wedge x \in B)$
- $x \in (A \cup B)$ if and only if $(x \in A \vee x \in B)$
- $x \in A - B$ if and only if $(x \in A \wedge x \notin B)$
- $x \in A \Delta B$ if and only if $(x \in A \wedge x \notin B) \vee (x \in B \wedge x \notin A)$
- $x \in A'$ if and only if $\neg(x \in A)$
- $X \in P(A)$ if and only if $X \subseteq A$

Venn Diagrams

Venn Diagrams are used to depict the various unions, subsets, complements, intersections etc. of sets.

Logic and Sets are closely related

Tautology

$p \vee q \leftrightarrow q \vee p$
 $p \wedge q \leftrightarrow q \wedge p$
 $p \vee (q \vee r) \leftrightarrow (p \vee q) \vee r$
 $p \wedge (q \wedge r) \leftrightarrow (p \wedge q) \wedge r$
 $p \vee (q \wedge r) \leftrightarrow (p \vee q) \wedge (p \vee r)$
 $p \wedge (q \vee r) \leftrightarrow (p \wedge q) \vee (p \wedge r)$
 $p \wedge \neg q \leftrightarrow p \wedge (\neg p \wedge q)$
 $p \wedge (\neg q \vee r) \leftrightarrow (p \wedge \neg q) \wedge (p \wedge r)$
 $p \wedge (\neg q \wedge r) \leftrightarrow (p \wedge \neg q) \vee (p \wedge r)$
 $p \wedge (q \wedge \neg r) \leftrightarrow (p \wedge q) \wedge (\neg p \wedge r)$
 $p \vee (q \wedge \neg r) \leftrightarrow (p \vee q) \wedge (\neg p \wedge r)$
 $p \wedge \neg (q \wedge r) \leftrightarrow (p \wedge \neg q) \vee (p \wedge r)$

Set Operation Identity

$A \cup B = B \cup A$
 $A \cap B = B \cap A$
 $A \cup (B \cap C) = (A \cup B) \cap C$
 $A \cap (B \cup C) = (A \cap B) \cup C$
 $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$
 $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
 $A - B = A - (A \cap B)$
 $A - (B \cap C) = (A - B) \cup (A - C)$
 $A - (B \cup C) = (A - B) \cap (A - C)$
 $A \cap (B - C) = (A \cap B) - (A \cap C)$
 $A \cup (B - C) = (A \cup B) - (C - A)$
 $A - (B - C) = (A - B) \cup (A \cap C)$

The above identities serve as the basis for an "algebra of sets".

Logic and Sets are closely related

Tautology

$p \wedge p \leftrightarrow p$
 $p \vee p \leftrightarrow p$
 $p \wedge (\neg q \wedge \neg q) \leftrightarrow p$
 $p \vee (\neg q \wedge \neg q) \leftrightarrow p$

Contradiction

$p \wedge \neg p$
 $p \wedge (q \wedge \neg q)$
 $p \wedge \neg p$

Set Operation Identity

$A \cap A = A$
 $A \cup A = A$
 $A - \emptyset = A$
 $A \cup \emptyset = A$

Set Operation Identity

$A - A = \emptyset$
 $A \cap \emptyset = \emptyset$
 $A - A = \emptyset$

The above identities serve as the basis for an "algebra of sets".

Theorem for Logic

Let p, q and r be propositions.

Idempotent laws:

$p \wedge p \equiv p$
 $p \vee p \equiv p$

Truth table:

p	$p \wedge p$	$p \vee p$
T	T	T
F	F	F

Theorem for Logic (cont.)

Double negation law:

$\neg \neg p \equiv p$

Commutative laws:

$p \wedge q \equiv q \wedge p$
 $p \vee q \equiv q \vee p$

Theorem for Logic (cont.)

Associative laws:

$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$
 $(p \vee q) \vee r \equiv p \vee (q \vee r)$

Distributive laws:

$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$
 $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

Absorption laws:

$p \wedge (p \vee q) \equiv p$
 $p \vee (p \wedge q) \equiv p$

Prove: Distributive Laws

p	q	r	$p \vee (q \wedge r)$	$(p \vee q) \wedge (p \vee r)$
T	T	T	T	T
T	T	F	T	T
T	F	T	T	T
T	F	F	T	T
F	T	T	T	T
F	T	F	F	F
F	F	T	F	F
F	F	F	F	F

Prove: Absorption Laws

p	q	$p \wedge (p \vee q)$	$p \vee (p \wedge q)$
T	T	T	T
T	F	T	T
F	T	F	F
F	F	F	F

Theorem for Logic (cont.)

De Morgan's laws:

$\neg(p \wedge q) \equiv (\neg p) \vee (\neg q)$
 $\neg(p \vee q) \equiv (\neg p) \wedge (\neg q)$

The **truth table** for $\neg(p \vee q) \equiv (\neg p) \wedge (\neg q)$

p	q	$\neg(p \vee q)$	$\neg p \wedge \neg q$
T	T	F	F
T	F	F	F
F	T	F	F
F	F	T	T

Exercise (6)

Given,

$$R = p \wedge (\neg q \vee r)$$

$$Q = p \vee (q \wedge \neg r)$$

State whether or not $R \equiv Q$.

Exercise (7)

Propositional functions p , q and r are defined as follows:

p is " $n = 7$ "
 q is " $a > 5$ "
 r is " $x = 0$ "

Write the following expressions in terms of p , q and r , and show that each pair of expressions is **logically equivalent**. State carefully which of the above laws are used at each stage.

(a) $((n = 7) \vee (a > 5)) (x = 0)$
 $((n = 7) (x = 0)) \vee ((a > 5) (x = 0))$

(b) $\neg((n = 7) (a \leq 5))$
 $(n \neq 7) \vee (a > 5)$

(c) $(n = 7) \vee (\neg((a \leq 5) (x = 0)))$
 $((n = 7) \vee (a > 5)) \vee (x \neq 0)$

Exercise (8)

Propositions p , q , r and s are defined as follows:

p is "I shall finish my Coursework Assignment"
 q is "I shall work for forty hours this week"
 r is "I shall pass Maths"
 s is "I like Maths"

Write each sentence in symbols:

(a) I shall not finish my Coursework Assignment.
 (b) I don't like Maths, but I shall finish my Coursework Assignment.
 (c) If I finish my Coursework Assignment, I shall pass Maths.
 (d) I shall pass Maths only if I work for forty hours this week and finish my Coursework Assignment.

Write each expression as a sensible (if untrue!) English sentence:

(e) $q \vee p$
 (f) $\neg p \Rightarrow \neg r$

Exercise (9)

For each pair of expressions, construct **truth tables** to see if the two compound propositions are logically equivalent:

(a) $p \vee (q \wedge \neg p)$
 $p \vee q$

(b) $(\neg p \wedge q) \vee (p \wedge \neg q)$
 $(\neg p \wedge \neg q) \vee (p \wedge q)$